



Algorithmic Problems in Free Groups

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Freely reduced words

A a finite alphabet $\neq \emptyset$
set

$$\bar{A} = \{ \bar{a} \mid a \in A \}$$

$$\bar{A} \cap A = \emptyset$$

$$\tilde{A} = A \cup \bar{A}$$

\tilde{A}^* = all words on alph \tilde{A}

ϵ is the empty word

Notation $\bar{a} = a$ for each $a \in A$

$$\begin{array}{ccc} x & \mapsto & \bar{x} \\ \tilde{A} & \longrightarrow & \tilde{A} \end{array} \text{ bij}$$

Want to see \bar{a} as an
(group) inverse of a

Let \sim be the congruence on \tilde{A}^*
generated by $a\bar{a} \sim \bar{a}a \sim 1$
for every $a \in \tilde{A}$

$$u, v \in \tilde{A}^*$$

$$u \sim v \text{ iff } \exists u_0 = u, u_1, \dots, u_n, u_n = v$$

$$\text{st. } u_i = pa\bar{a}q \text{ and } u_{i+1} = pq \\ \text{or } u_i = pq \text{ and } u_{i+1} = pa\bar{a}q$$

Exple $a b \bar{a} a \bar{b} a$
 $\sim a b \bar{b} \bar{a}$
 $\sim \boxed{abc}$ reduced
 $\sim a \bar{a} a b a$

A word is reduced if it contains
 (freely) no $a\bar{a}$ or $\bar{a}a$
 ($a \in \hat{A}$)

Prop Every word $u \in \hat{A}^*$ is
 \sim -equivalent to a
unique reduced word
written

$u = p a \bar{a} a q$ red(u)
 \downarrow
 $p a q$

B an alphabet, elements are called letters

words are finite sequences of letters

sequences of length 1 = B

sequence of length 0 = the empty word

B^* = all words

$1 \in B^*$, $1 \notin B$

Free group

$F(A)$ = the set of all reduced words in A^*

$$u, v \in F(A)$$

$$u \cdot v = \text{red}(uv)$$

This operation is associative

• \sim is a congruence

that is: if $u_1 \sim u'_1$

$$u_2 \sim u'_2$$

then $u_1 u_2 \sim u'_1 u'_2$

• $\text{red}(u)$ is unique

1 is an identity element

$$(1 \cdot u = u \cdot 1 = u)$$

$$a \cdot \bar{a} = \bar{a} \cdot a = 1 \quad \text{for } a \in A^*$$

$$\bar{a} = a^{-1}$$

$$u = a_1 \dots a_n$$

$$a_i \in A^*$$

$$u^{-1} = a_n^{-1} a_{n-1}^{-1} \dots a_1^{-1}$$

$$= \bar{a}_n \bar{a}_{n-1} \dots \bar{a}_1$$

$F(A)$ is a group
called the free group on A

$F(A) = A^{\ast} / \sim$ quotient monoid, which turns out to be a gp.

Let G be any group and $\varphi: A \rightarrow G$
be any map
Then φ extends uniquely to a
group homomorphism $\hat{\varphi}: F(A) \rightarrow G$

$$a b a c \xrightarrow{\hat{\varphi}} \varphi(a) \varphi(b) \varphi(a) \varphi(c)$$

$$\hat{\varphi}(\bar{a}) = \varphi(a)^{-1}$$

$a \cdot \bar{a}$ in $F(A)$

is 1 so $\hat{\varphi}(a \cdot \bar{a}) = \hat{\varphi}(a) \hat{\varphi}(\bar{a}) = 1$

$$a \bar{a} \mapsto \varphi(a) \varphi(\bar{a})$$

Make that $\varphi(\bar{a}) = \varphi(a)^{-1}$

What if I change alphabet

$F(A)$ isomorphic $F(B)$

?

Theorem $F(A) \cong F(B) \iff |A| = |B|$

$\text{card}(A) = \text{card}(B)$

Sketch of proof of \Rightarrow

$$F(A) \xrightarrow{\tau_A} \left(\frac{\mathbb{Z}}{2\mathbb{Z}}\right)^A = \bigoplus_{a \in A} \frac{\mathbb{Z}}{2\mathbb{Z}} a$$

$$a \longmapsto 1 \cdot a$$

$a \in A$

vector space over the field $\frac{\mathbb{Z}}{2\mathbb{Z}}$ with basis A

$|A| = \text{rank of } F(A)$

τ_A is a surjective group homomorphism

if $F(A) \xrightarrow{\varphi} F(B)$ is an isomorphism

$$\begin{array}{ccc} F(A) & \xrightarrow{\varphi} & F(B) \\ \tau_A \downarrow & & \downarrow \tau_B \\ V_A \left(\frac{\mathbb{Z}}{2\mathbb{Z}}\right)^A & \xrightarrow{\varphi_2} & \left(\frac{\mathbb{Z}}{2\mathbb{Z}}\right)^B V_B \end{array}$$

φ_2 isom

$|A| = \dim V_A$
 $|B| = \dim V_B$

$F(A)$ has rank $|A|$

A is a basis of $F(A)$

by defⁿ: every el['] of $F(A)$
can be written in a unique
way as a reduced word
on \tilde{A} (in \tilde{A}^{*0})

$$A = \{a, b\}$$

$\{a, b^{-1}\}$ is also a basis

$\{ab, b\}$ _____

~~$\{b^{-1}ab, b\}$~~ _____

$\{\underline{b^{-1}ab^2}, \underline{b^{-1}ab}\}$ _____

if $\{u, v\}$ is a basis, then

$\{u, v^{-1}\}$ is a basis

$\{uv, v\}$ is a basis

All bases have cardinality $\text{rank}(F(A))$
if $u_i \rightarrow u_{i+1}$ is a basis of $F(A)$ $= |A|$

If $\{u_1, \dots, u_n\}$ is a basis of $F(A)$

and $B = \{b_1, \dots, b_n\}$ is a set with n elt

then $F(B) \rightarrow F(A)$ is an
 $b_i \mapsto u_i$ isomorphism

Problem given u_1, \dots, u_r in $F(A)$
(where $l(A) = r$)
decide whether $\{u_1, \dots, u_r\}$ is a basis
of $F(A)$.

$$A = \langle a, b \rangle$$

$\{aba, ba^{-1}b\}$ basis?

Subgroups of free groups

Theorem Every (finitely generated) subgroup of $F(A)$ is free
(Nielsen)
early 20th
century)

Contrary to vector spaces

if H is a subgroup of $F(A)$
 $\text{rank}(H)$ may be greater than $|A|$
may even be ∞

Example in $F(a, b)$

any subset $\{b^i a b^{-i} \mid i \in \mathbb{Z}\}$ is a basis
of the subgroup it generates

Problems given $\underline{g_1, g_2, \dots, g_n}$ in $F(A)$

let $H = \langle g_1, \dots, g_n \rangle$

1) is $\{g_1, \dots, g_n\}$ a basis of H ?

2) $\text{rk}(H)$?

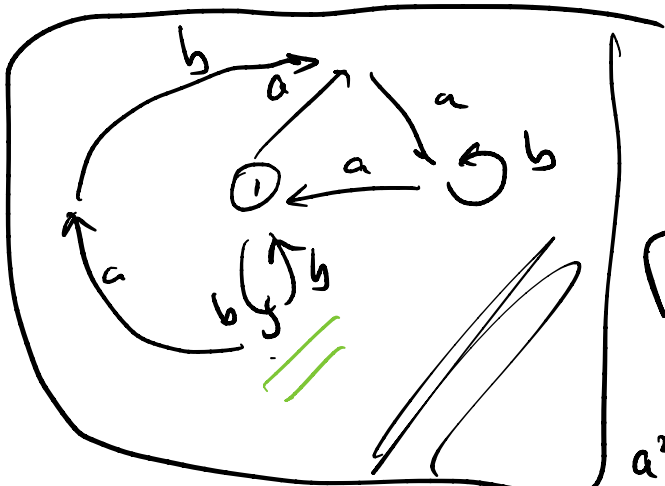
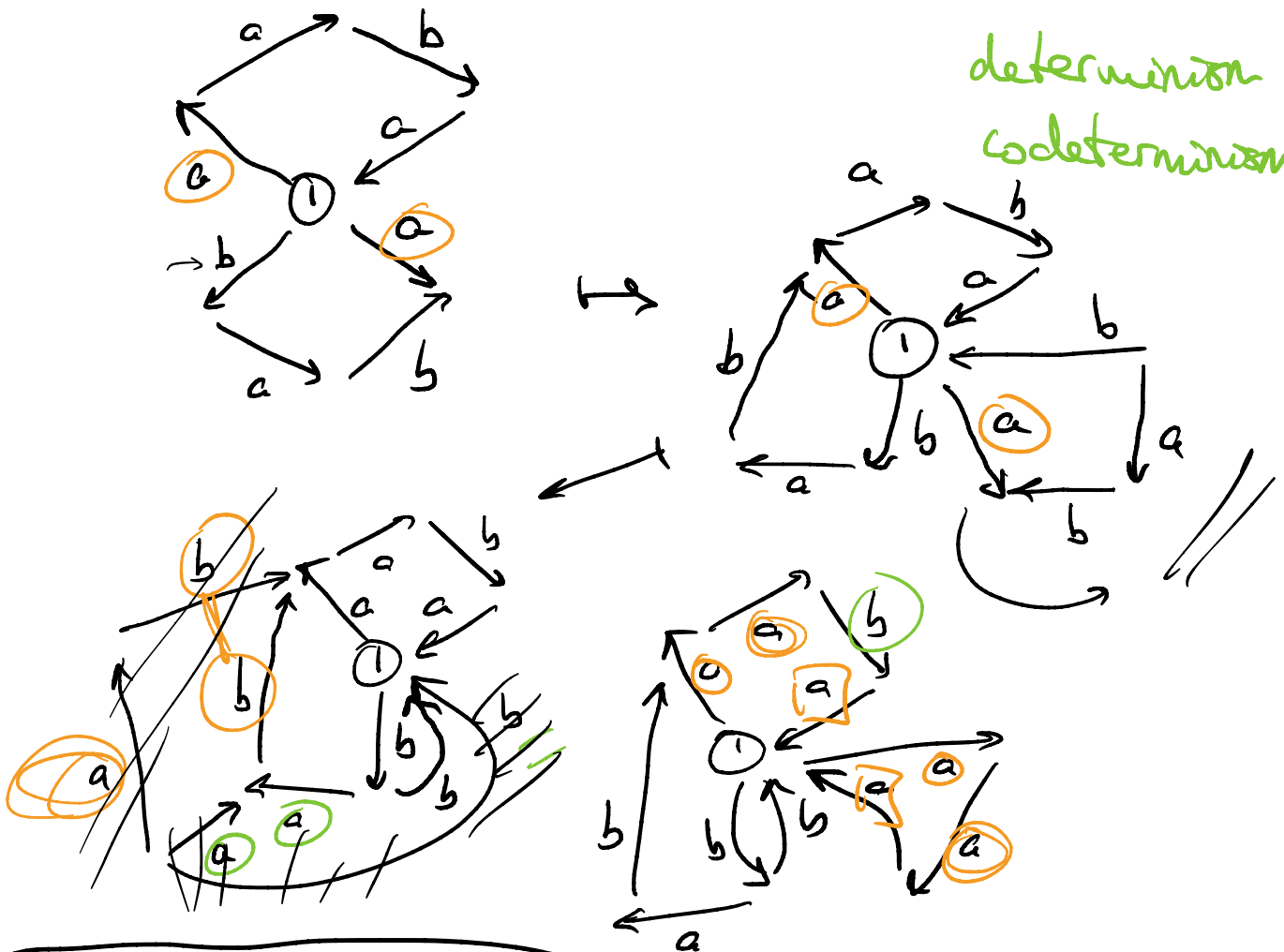
3) is $g \in H$?

Stallings graph of a subgroup

$H \longmapsto \Gamma(H)$ a finite connected subgp graph

$H = \langle a^2ba, babab^{-1}, \underline{b^{-1}aba^{-1}}, a^3, \underbrace{b^2} \rangle$

determination
codetermination



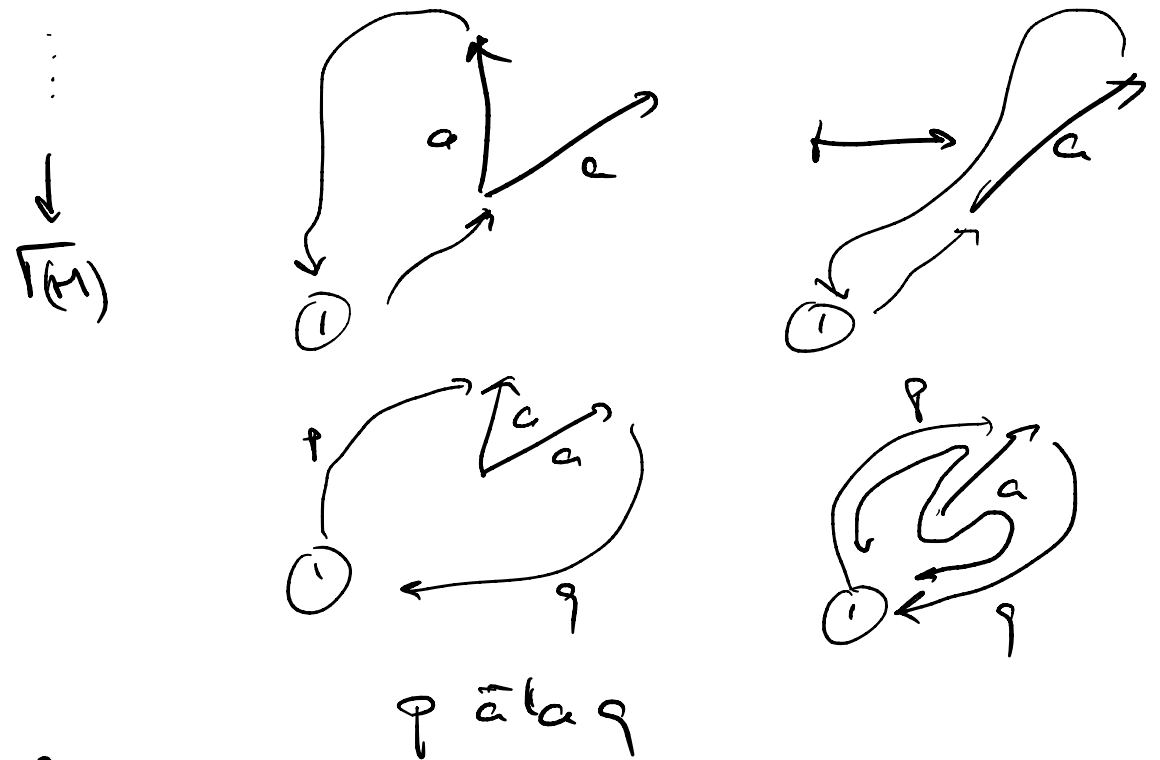
$a^2b \notin H$
 $bababa \in H$

$\Gamma(H)$ does not depend on the set of generators of H
It depends on H only

$\Gamma_0 =$ bouquet of circles

\downarrow folding $L(\Gamma_1) =$ language in A^*

\downarrow folding $\forall u \in H \exists v \in L(\Gamma_0)$ such that $u = \text{red}(v)$



Thm A reduced word is in H

iff it is accepted by $\Gamma(H)$ (seen as an automaton)

You start with generators of H
 (h_1, \dots, h_n)

construct $\Gamma(H)$

given $g, h_1, \dots, h_n, H = \langle h_1, \dots, h_n \rangle$
 $\Rightarrow g \in H?$

Prop $g \in H$ iff g can be read in $\Gamma(H)$
as a circuit at the base
vertex

$\{g_1, g_2, \dots, g_r\}$ a basis of $F(A)$
(where $|A| = r$)

Let $H = \langle g_1, \dots, g_r \rangle$

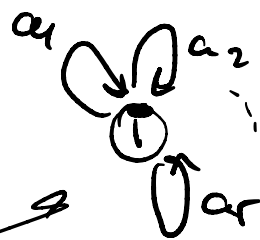
$\Gamma(H)$

H is $F(A) \Leftrightarrow$

$\Gamma(H) = \Gamma(F(A))$

$\{g_1, \dots, g_r\}$ is a
basis of $F(A)$

$\Rightarrow \Gamma(H) =$



Computation of \cap -sections

H, K 2 subgrps of $F(A)$

Compute $H \cap K$

The St. graph for $H \cap K$
is obtained from $\Gamma(H)$
 $\rightarrow \Gamma(K)$ by the same
product construction used
to recognize $L_1 \cap L_2$
when L_1, L_2 are regular
languages

Shallberg
1983

Cardan: if $rk(H), rk(K) < \infty$
then $rk(H \cap K) < \infty$
(Kasson's theorem 1954)

Surveys

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see also my webpage