#### Automorphism Group of Graphs

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# Graphs

- G=(V,E); V=set of vertices; E=set of edges
- order of G = V = n



Figure: V={a,b,c,d,e,f} E={ab,ac,ad,be,cf,bf}

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- A permutation of the set  $\Omega$  is a bijective mapping  $g:\Omega\to\Omega$
- The composition  $g_1g_2$  of two permutations  $g_1$  and  $g_2$  is

$$v(g_1g_2)=(vg_1)g_2$$
 for each  $v\in \Omega$ 

- G is closed under composition: if  $g_1, g_2 \in G$  then  $g_1g_2 \in G$ ;
- G contains the *identity* permutation 1, defined by v1 = v for  $v \in \Omega$ ;
- G is closed under inversion, where the inverse of g is the permutation  $g^{-1}$  defined by the rule that  $vg^{-1} = w$  if wg = v;

#### Definition

An automorphism on a graph G is a bijection  $\phi : V(G) \to V(G)$  such that  $uv \in E(G)$  if and only if  $\phi(u)\phi(v) \in E(G)$ .

Note that graph automorphisms preserve *adjacency*. In layman terms, a graph automorphism is a symmetry of the graph.

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Figure: Graph G

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One automorphism switches vertices a and c.



Figure:  $\alpha = (a, c)(b)(d)$ 

A D N A B N A B N A B N

One automorphism switches vertices b and d.



Figure:  $\beta = (a)(b, d)(c)$ 

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The final automorphism switches vertices a and c and also switches b and d.



Figure:  $\alpha\beta = (a, c)(b, d)$ 

#### Definition

Let Aut(G) denote the set of all automorphisms on a graph G. Note that this forms a group under function composition.

- Aut(G) is closed under function composition.
- Function composition is associative on Aut(G).
- There is an identity element in Aut(G). This is mapping e(v) = v for all v ∈ V(G).
- For every σ ∈ Aut(G), there is an inverse element σ<sup>-1</sup> ∈ Aut(G).
  Since σ is a bijection, it has an inverse. By definition, this is an automorphism.

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- Graph automorphisms are degree preserving. In other words, for all  $u \in V(G)$  and for all  $\phi \in Aut(G)$ ,  $deg(u) = deg(\phi(u))$ .
- Graph automorphisms are distance preserving. In other words, for all  $u, v \in V(G)$  and for all  $\phi \in Aut(G)$ ,  $d(u, v) = d(\phi(u), \phi(v))$ .
- The automorphism group of G is equal to the automorphism group of the complement  $\overline{G}$ .

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## Dihedral Group

The 5-cycle  $C_5$  has ten automorphisms, realized geometrically as the rotations and reflections of a regular pentagon.



#### Definition

There exists a function 'f' from vertices of  $G_1$  to vertices of  $G_2$ [ $f: V(G_1) \rightarrow V(G_2)$ ], such that (i) f is a bijection (both one-one and onto) (ii) f preserves adjacency of vertices, i.e., if the edge  $\{u, v\} \in G_1$ , then the edge  $\{f(u), f(v)\} \in G_2$ , then  $G_1 \equiv G_2$ .

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Figure: *f* is an isomorphism

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#### Question 1

**Graph isomorphism Instance:** Graphs G and H **Question:** Is  $G \cong H$ ?

#### Question 2

Automorphism group Instance: A graph GOutput: generating permutations for Aut(G).

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#### Reduction

#### Theorem

With **Graph-Iso** as an oracle, there is a polynomial time algorithm for **Graph-Aut** and vice-versa.

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- given two graphs  $G_1$  and  $G_2$
- Create  $G = G_1 \cup G_2$
- $G_1$  and  $G_2$  are connected
- if any of the generators of Aut(G) interchanged a vertex in  $G_1$  with one in  $G_2$ , then connnectivity should force  $G_1 \cong G_2$
- if not  $G_1$  and  $G_2$  connected
- $G_1 \cong G_2 \iff \overline{G_1} \cong \overline{G_2}$
- either  $G_1$  or  $\overline{G_1}$  has to be connected

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- Constructing the automorphism group is at least as difficult (in terms of its computational complexity) as solving the graph isomorphism problem.
- Similar to the graph isomorphism problem, it is unknown whether it has a polynomial time algorithm or it is NP-complete.

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# Questions...?



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