

PLENITUDE

SREEHARI K

JOINT WORK WITH KAMAL LODAYA

IISc, BANGALORE



Reasoning with *Few* and *Many*

- *Many* and *Few* are indefinite plural descriptions.
- For instance, by “Many apples in the refrigerator” one might mean there are a dozen, whereas by “Many apples in the orchard” one might mean there are 100.
- There are various accounts of reasoning with *Many* in literature. For instance, using syllogisms, generalized quantifiers, etc.
- Our approach is propositional, i.e., we write ‘Many(**a**)’ to indicate that the plural term ‘**a**’ are many. Similarly for *Few*.
- After defining *Few* and *Many*, we prove the completeness theorem for the extended system.

Mid-Plural Logic - Syntax

Plural logic is an extension of first-order logic, in which terms can be of either singular (denoting *an* individual) or plural (denoting individuals) sort.

Terms

- Singular (x, y, z, \dots) and Plural ($\mathbf{x}, \mathbf{y}, \mathbf{z}, \dots$) Variables
- Constants (c, \dots)
- $f a_1, a_2, \dots, a_n$ – where f is an n -ary function symbol and a_1, a_2, \dots, a_n are terms
- $x:A$ – where x is a variable and A is a formula. This is called exhaustive description
- Plural union (\mathbf{a}, \mathbf{b}), plural intersection ($\mathbf{a}.\mathbf{b}$), and plural complement ($-\mathbf{a}$).

Mid-Plural Logic – Syntax (contd.)

Formulas

- Pa_1, a_2, \dots, a_n – where P is an n -ary predicate symbol and a_1, a_2, \dots, a_n are terms
- $a \preceq b$ – where a and b are terms. Read as a is among/included in b .
- $A \supset B, \neg A, \dots$
- $\exists xA, \forall xA$

Mid-Plural Logic – Syntax (contd.)

Some Definitions

- $E(\mathbf{a}) =^{\text{def}} \exists x (x \leq \mathbf{a})$ – Says that the term \mathbf{a} denotes.
- $S(\mathbf{a}) =^{\text{def}} \exists x (x = \mathbf{a})$ – Says that the term \mathbf{a} is singular
- $true =^{\text{def}} (x = x)$ – Thus $x: true$ denotes all the individuals
- One can also define the terms plural union (\mathbf{a}, \mathbf{b}) , plural intersection $(\mathbf{a}.\mathbf{b})$, and plural complement $(-\mathbf{a})$ using inclusion and exhaustive description.

Mid-Plural Logic – Axiom System

Oliver and Smiley provides a Hilbert-style axiom system. Additional axioms are needed for $:$ and \leq . These are –

- $\forall y(y \leq x : A(x) \equiv A(y))$ - where $A(y)$ has free y wherever $A(x)$ has free x .
- $\mathbf{a} \leq \mathbf{b} \equiv \exists \mathbf{a} \wedge \forall x(x \leq \mathbf{a} \supset x \leq \mathbf{b})$ - where x is not free in \mathbf{a} or \mathbf{b} .
- Sx

Note: One cannot quantify over plurals without sacrificing axiomatizability. This is because allowing quantification over plurals enables one to define a categorical, finitely axiomatized version of Peano Arithmetic.

Mid-Plural Logic – Semantics

- We have individuals (domain)
- A valuation val relating objects to terms, relations to predicates, and functions to function symbols, as follows
 - $val\ x$ is an individual and $val\ \mathbf{x}$ are individuals.
 - $val\ P$ is an n -place relation among individuals, where P is an n -ary predicate.
 - $val\ f$ is an n -place function on individuals, where f is an n -ary function symbol.
 - $val\ x:A$ are the x -variants $val'\ x$ for every x -variant val' satisfying A .

The rules for satisfaction are similar to first-order logic.

For inclusion, $val \models \mathbf{a} \leq \mathbf{b}$ iff each individual in $val\ \mathbf{a}$ is in $val\ \mathbf{b}$ as well.

Introducing *Few* and *Many*

We add three new unary predicate symbols, '*Many*', '*Few*', and ' $Bet_{Few,Many}$ ' to the signature.

$Bet_{Few,Many}$ represents the *penumbra*, i.e., areas of indeterminacy where one cannot commit to either *Few* or *Many*.

Therefore, for each term **a**, we have (new) formulas –

- *Many* (**a**), which says that **a** are many.
- *Few* (**a**), which says that **a** are few.
- $Bet_{Few,Many}$ (**a**), which says that **a** are more than few, but not many.

Example Sentences

1. “There are many apples in the refrigerator” – $Many(x:(apple(x) \wedge in.fridge(x)))$
2. “Every actor has many talents” – $\forall x (Actor(x) \supset Many(\mathbf{talents}(x)))$

Note that in (1), “apple” and “in.fridge” are singular predicates.

In (2), **talents** is a multi-valued function. “Actor” is a singular predicate.

Extended Axiom System

To Oliver and Smiley's Hilbert-style system, we add the following axioms for the new symbols.

1. $\text{Many } (x:\text{true})$ – Says that the domain has *many* individuals.
2. $\text{Sa} \supset \text{Few } \mathbf{a}$ – Singletons are few.
3. $\mathbf{a} \leq \mathbf{b} \supset (\text{Many } \mathbf{a} \supset \text{Many } \mathbf{b}) \wedge (\text{Few } \mathbf{b} \supset \text{Few } \mathbf{a})$ – Expresses Anti-monotonicity of *few* and Monotonicity of *many*.

From (3), we can derive $\text{Many } \mathbf{a.b} \supset \text{Many } \mathbf{a}$ and $\text{Many } \mathbf{a} \supset \text{Many } \mathbf{a,b}$

Extended Axiom System (contd.)

4. $Few\ a \supset \neg\ Many\ a$ – *Few* is a degree less than *Many*. Hence partitions plurals into three degrees, which are the new predicates.
5. $\neg\ Many\ a \supset Many\ -a$ – Relates negation to *Many*, as well as partitioning *Many* into three degrees, i.e., $Many\ a \wedge Many\ -a, \dots$
6. (i) $\forall x (Few\ a \supset (Few\ a,x \vee Bet_{Few,Many}\ a,x))$
(ii) $\forall x (Bet_{Few,Many}\ a \supset ((Bet_{Few,Many}\ a,x \vee Many\ a,x) \wedge (Bet_{Few,Many}\ a-x \vee Few\ a-x)))$
(iii) $\forall x (Many\ a \supset (Bet_{Few,Many}\ a-x \vee Many\ a-x))$
– Sets up degree ordering in the presence of penumbral predicate.
 $Few < Bet_{Few,Many} < Many$

Semantics for the Extension

We provide an intuitive Tarskian semantics for the three degrees introduced.

- $val \models \textit{Many } \mathbf{a}$ iff there are many among $val(\mathbf{a})$
- $val \models \textit{Few } \mathbf{a}$ iff there are only few among $val(\mathbf{a})$
- $val \models \textit{Bet}_{\textit{Few}, \textit{Many}} \mathbf{a}$ iff there are more than few but less than many among $val(\mathbf{a})$

This is justified by noting that *Few*, *Many* sets up comparison classes.

Thus the first statement says that *val* satisfies *Many* \mathbf{a} iff there are more than $m_{\mathbf{a}}$ individuals in $val(\mathbf{a})$, where $m_{\mathbf{a}}$ is indefinite.

Completeness Theorem

- Oliver and Smiley prove the completeness theorem for mid plural logic. They follow Henkin's method of using Lindenbaum construction.
- To prove the contrapositive of completeness theorem, i.e., if $val \not\models C$ then $val \not\models C$, one starts by adding countable henkin constants h to the signature.
- Then, using an enumeration of formulas, construct a maximally consistent set Δ such that $\Delta \not\models C$ (for each existential formula for a fresh witness h add $E(h), S(h)$).
- Use Δ to construct a model with domain consisting of the henkin constants as individuals.

Key Lemma : For all terms \mathbf{a} and formulas A , $val \mathbf{a}$ are (weakly) identical to \mathbf{a}^* (denoted by h s.t. $h \leq \mathbf{a}$ is in Δ) and val satisfies A iff A is one of Δ .

Lindenbaum Construction for the Extended System

When one of $Many\ a$, $Bet_{Few, Many}\ a$, $Few\ a$ is encountered in Δ , do the following.

Many a

- Find maximal \mathbf{b} such that $\Delta \vdash \mathbf{b} \leq \mathbf{a} \wedge \neg Many\ \mathbf{b}$
- Add sufficiently many *distinct* $h \leq \mathbf{a}$, $h \leq -\mathbf{b}$, $\exists x(x = h)$ to Δ if required, to justify $Many\ a$ (Plenitude conditions).

$Bet_{Few, Many}\ a$

- Find maximal \mathbf{b} and minimal \mathbf{c} such that $\Delta \vdash \mathbf{b} \leq \mathbf{a} \wedge Few\ \mathbf{b}$, $\Delta \vdash \mathbf{a} \leq \mathbf{c} \wedge Many\ \mathbf{c}$
- If required, add henkin constants and preserve separation of counts (as before).

Few a

- Find minimal \mathbf{b} such that $\Delta \vdash \mathbf{b} \leq \mathbf{a} \wedge \neg Few\ \mathbf{b}$
- Preserve separation of counts, and the count must be below distinct $h \leq \mathbf{b}$.

Proof Sketch of Key Lemma for *Many*

Statement :

val satisfies *Many a* if and only if the formula *Many a* is one of the truth set Δ

Proof Sketch :

- If *Many a* is in Δ , by our construction and the plenitude condition, there are many *distinct* henkin constants such that $h \leq a$.
- For the converse, if \neg *Many a* is in Δ , the construction and the partition axiom ensures that the *distinct* henkin constants such that $h \leq a$ are not many.

Further Work

- Consider the formulas $\forall x (A(x, \mathbf{x}) \wedge \text{Many } \mathbf{x})$ and $(\forall x A(x, \mathbf{x}) \wedge \text{Many } \mathbf{x})$. We might expect the comparison count of \mathbf{x} for *Many* to depend on the values of x in the first case but not the second. But they are equivalent in our system. Can one remedy this, for instance, by employing Skolem functions?
- Incorporating *Most*.
- *Few* and *Many* are vague predicates. Problems of discerning boundary pertain to these predicates as well. We have provided an intuitive Tarskian semantics here. How does one employ this idea of *indefinite descriptions* in the case of, say, the vagueness of Orange-Red?
- Extending the idea of indefinite description to finite (>3) and countably infinite degrees.

THANK YOU
